## LECTURE 6: SIMPLEX ALGORITHM (II)

**Today:** Examples of the Simplex Algorithm, which is an efficient way of finding the optimal vertex.

## 1. SIMPLEX ALGORITHM EXAMPLE

Example 1:		
	$\max z = 2x_1 + 5x_2$	
	subject to $2x_1 - x_2 \le 4$	(1)
	$x_1 + 2x_2 \le 9$	2
	$-x_1 + x_2 \le 3$	3
	$x_1 \ge 0$	4
	$x_2 \ge 0$	5

**Optional Picture:** See next page

**STEP 1:** Start at (0, 0)

Current Vertex:  $\{(4), 5\}$  (the constraints satisfied with equality)

## **Objective Value:** 0

Not optimal! we can increase  $x_1$  or  $x_2$  to increase the value of z. We choose  $x_2$  since 5 is the biggest number (= fastest increase)

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Move: Increase  $x_2$  until you hit a constraint

Question: Which constraint (1), (2), (3) is hit first?

We are moving along the line  $(0, x_2)$ , so let's set  $x_1 = 0$  in each constraint and see which  $x_2$  value is obtained:

(1)  $2(0) - x_2 = 4 \Rightarrow x_2 = -4 \times$ (2)  $0 + 2x_2 = 9 \Rightarrow x_2 = \frac{9}{2} = 4.5$ (3)  $-0 + x_2 = 3 \Rightarrow x_2 = 3$ 

(3) is hit first since it gives the smallest hitting time  $x_2 = 3$ 

So you hit the new vertex  $\{(4), (3)\} = (0, 3)$ 

Note: What helped us in this problem is that we started at the origin and our constraints were  $x_1 \ge 0$  and  $x_2 \ge 0$ .

**Idea:** Introduce new coordinates  $(y_1, y_2)$  so that (0,3) becomes the new origin and the constraints (4) and (3) become  $y_1 \ge 0$  and  $y_2 \ge 0$ 

#### Fact:

If the constraint is  $a^T x \leq b$ , then  $y = b - a^T x$ 

(Think  $b - a^T x \ge 0$ , making the variables positive, like slack variables)

$$\begin{array}{ll}
(4) & x_1 \ge 0 \Rightarrow y_1 = x_1 \\
(3) & -x_1 + x_2 \le 3 \Rightarrow 3 + x_1 - x_2 \ge 0 \Rightarrow y_2 = 3 + x_1 - x_2
\end{array}$$

**Coordinates:** 

$$\begin{cases} y_1 = x_1 \\ y_2 = 3 + x_1 - x_2 \end{cases}$$

**Geometric Interpretation:** Here  $y_1$  and  $y_2$  are the directions perpendicular to the constraints (4) and (3), as in the following picture:



# **Rewrite LP in terms of** $y_1$ and $y_2$ :

$$\begin{cases} x_1 = y_1 \\ y_2 = 3 + x_1 - x_2 \Rightarrow x_2 = 3 + x_1 - y_2 = 3 + y_1 - y_2 \\ z = 2x_1 + 5x_2 = 2y_1 + 5(3 + y_1 - y_2) = 15 + 7y_1 - 5y_2 \end{cases}$$

And similarly, you rewrite the constraints, and so the LP becomes

$\max z = 15 + 7y_1 - 5y_2$	
Subject to $y_1 + y_2 \le 7$	1
$3y_1 - 2y_2 \le 3$	2
$y_2 \ge 0$	3
$y_1 \ge 0$	4
$-y_1 + y_2 \le 3$	5

**STEP 2:** Repeat, but for the new vertex

Current Vertex:  $\{(4), (3)\}$ 

**Objective Value:** z = 15

Not optimal, because of the 7 in  $7y_1$ 

**Move:** Increase  $y_1$  (so ④ is released)

If  $y_2 = 0$  then

(1)	$y_1 + 0 = 7 \Rightarrow y_1 = 7$
2	$3y_1 - 0 = 3 \Rightarrow y_1 = 1$
5	$-y_1 + 0 = 3 \Rightarrow y_1 = -3 \times$

So we hit the constraint ② first

**Coordinates:** Vertex  $\{(2), (3)\} = (1, 0)$  has coordinates  $(z_1, z_2)$ , where:

$$\begin{cases} z_1 = 3 - (3y_1 - 2y_2) = 3 - 3y_1 + 2y_2 \\ z_2 = y_2 \end{cases}$$

## **Rewrite LP:**

$$y_{2} = z_{2}$$

$$z_{1} = 3 - 3y_{1} + 2y_{2} \Rightarrow 3y_{1} = 3 - z_{1} + 2y_{2} = 3 - z_{1} + 2z_{2} \Rightarrow y_{1} = 1 - \frac{1}{3}z_{1} + \frac{2}{3}z_{2}$$

$$\begin{cases} y_{1} = 1 - \frac{1}{3}z_{1} + \frac{2}{3}z_{2} \\ y_{2} = z_{2} \end{cases}$$

Then in terms of our new variables, the LP problem becomes



Note: Strictly speaking, it's not necessary here to rewrite the constraints here in terms of  $z_1$  and  $z_2$  because we'll see below that z = 22is optimal.

#### **STEP 3:**

Current Vertex:  $\{2,3\}$ 

### **Objective Value:** 22

This is optimal because all the coefficients of  $z_1$  and  $z_2$  are negative!

That is, no matter which direction you move in, you will only *decrease* the value of z (remember  $z_1 \ge 0$  and  $z_2 \ge 0$ )

And so the maximal value is 22

**Optimal Vertex:** The optimal vertex is (0,0) in z-coordinates

Now use our formulas for  $y_1, y_2$  and  $x_1, x_2$  to write it back in x-coordinates

$$\begin{cases} z_1 = 0\\ z_2 = 0 \end{cases}$$

$$\begin{cases} y_1 = 1 - \frac{1}{3}z_1 + \frac{2}{3}z_2 = 1 - 0 + 0 = 1\\ y_2 = z_2 = 0 \end{cases}$$

$$\begin{cases} x_1 = y_1 = 1\\ x_2 = 3 + y_1 - y_2 = 3 + 1 - 0 = 4 \end{cases}$$

So our optimal vertex is  $(x_1, x_2) = (1, 4)$  and our optimal value is z = 22.

#### 2. HIGHER DIMENSIONAL EXAMPLE

This method is so powerful that we can apply this to 3+ variables



(Will only do **STEP 1**, just to show you how it works)

**STEP 1:** Start at (0, 0, 0, 0, 0)

Current Vertex:  $\{(4), (5), (6), (7), (8)\}$ 

## **Objective Value:** z = 0

Not optimal, because of the 3, 4, and 4 coefficients

We can increase either  $x_2$  or  $x_5$ , let's do it for  $x_2$ 

**Move:** Increase  $x_2$  (so (5) is released)

If  $x_1, x_3, x_4$ , and  $x_5$  are 0 then

(2) 
$$2x_2 + 0 = 12 \Rightarrow x_2 = 6$$
  
(3)  $3(0) + 2x_2 + 0 = 18 \Rightarrow x_2 = 9$ 

So we hit the constraint (2) first

**Coordinates:** Vertex  $\{(2), (4), (6), (7), (8)\}$  has coordinates  $(y_1, y_2, y_3, y_4, y_5)$ 

$$\begin{cases} y_1 = x_1 \\ y_2 = 12 - 2x_2 - x_4 \\ y_3 = x_3 \\ y_4 = x_4 \\ y_5 = x_5 \end{cases}$$

**Rewrite LP:** 

$$\begin{cases} x_1 = y_1 \\ 2x_2 = 12 - y_2 - x_4 \Rightarrow x_2 = 6 - \frac{1}{2}y_2 - \frac{1}{2}y_4 \\ x_3 = y_3 \\ x_4 = y_4 \\ x_5 = y_5 \end{cases}$$

$$z = 3x_1 + 4x_2 + 5x_5 = 3y_1 + 4\left(6 - \frac{1}{2}y_2 - \frac{1}{2}y_4\right) + 5y_5 = 24 + 3y_1 - 2y_2 - 2y_4 + 4y_5$$

And similarly for the constraints

Then in terms of our new variables, the LP problem becomes

$\max z = 24 + 3y_1 - 2y_2 - 2y_4 + 4y_5$	
subject to $y_1 + y_3 \le 4$	(1)
$y_2 \ge 0$	2
$3y_1 - y_2 - y_4 + y_5 \le 6$	3
$y_1 \ge 0$	4
$6 - \frac{1}{2}y_2 - \frac{1}{2}y_4 \ge 0$	5
$y_3,y_4,y_5 \geq 0$	(6) - (8)

**STEP 2:** Then you would continue, increase  $y_1$  because this gives you a positive coefficient, then release ③ until you hit constraint ③ with  $y_1 = 2$ , change coordinates, check if this is optimal, etc.

Note: Had you started by increasing  $x_5$ , you would have immediately obtained the optimal vertex (0, 0, 0, 0, 18) with optimal value z = 72. You can actually guess  $x_5$  beforehand because it would give you a bigger hitting time  $x_5 = 18$  compared to  $x_2 = 6$ .

#### 3. MATLAB IMPLEMENTATION

The cool thing is that you won't ever have to perform the simplex algorithm by hand (except for the homework and the exam) because there is a package on MATLAB that does it for you!

Example 3:

 $\max z = 2x_1 + 5x_2$ subject to  $2x_1 - x_2 \le 4$  $x_1 + 2x_2 \le 9$  $-x_1 + x_2 \le 3$  $x_1 \ge 0$  $x_2 \ge 0$ 

In order to input this into MATLAB, we need to write this in the form z = 0 and  $Ax \leq b$ :

$$\max z - 2x_1 - 5x_2 = 0$$
  
subject to 
$$2x_1 - x_2 \le 4$$
$$x_1 + 2x_2 \le 9$$
$$-x_1 + x_2 \le 3$$
$$-x_1 \le 0$$
$$-x_2 \le 0$$

 $f = \begin{bmatrix} -2 & -5 \end{bmatrix}$  Coefficients of z, row vector

$$A = \begin{bmatrix} 2 & -1 \\ 1 & 2 \\ -1 & 1 \\ -1 & 0 \\ 0 & -1 \end{bmatrix}$$

$$b = \begin{bmatrix} 4\\9\\3\\0\\0 \end{bmatrix}$$
 Column vector

f = [-2,-5]
A = [2 -1;1 2; -1 1; -1 0; 0 -1]
b = [4 9 3 0 0]'
x = linprog(f,A,b)

Then gives you x = (1, 4), as we found above. If you want the optimal value, you just type

$$[x,z] = linprog(f,A,b)$$

This gives you x = (1, 4) and z = -22, and then the optimal value is z = -(-22) = 22

Note: We have to put that minus-sign because MATLAB solves min z instead of max z. The two are related via the equation

$$\max z = -\min(-z)$$

This is also why we have to input -2 and -5 in f instead of 2 and 5

Note: Similarly, if you apply this to the previous example, you would get x = (0, 0, 0, 0, 18) and z = -72, which means the max is 72.